include and the

nt main(){

std::cout <<

std::vector< std::iota(my)

# Concurrency

std::cout << ry binder of auto 1: ry vec. std::cout

std::cout << "\n\n";

Rainer Grimm Training, Mentoring, and Technology Consulting

#### Definition

"Each pattern is a three-part rule, which expresses a relation between a certain context, a problem, and a solution." (Christopher Alexander)

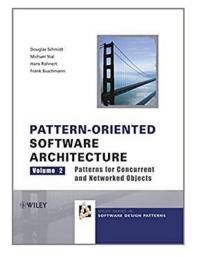
# Three Types of Patterns

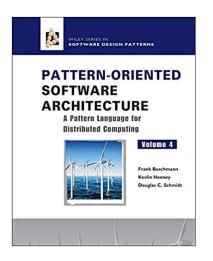
- Architecture pattern
  - Fundamental structure
  - Software system
- Design pattern
  - Interplay of components
  - Focus on a subsystem
- Idiome
  - Implementation of an architecture or design pattern in a programming language.

#### Components of a Pattern

- 1. Name
- 2. Also known as
- 3. Summary
- 4. Motivation
- 5. Context
- 6. Interaction
- 7. Solution
- 8. Example
- 9. Consequenses
- 10. Related pattern
- 11. Known usages

#### Pattern-Oriented Software Architecture (Volume 2 and 4)





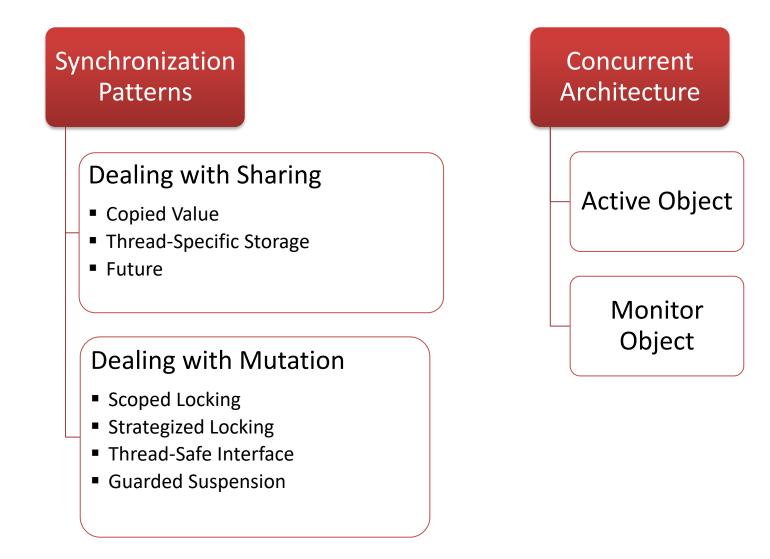
Doug Lea

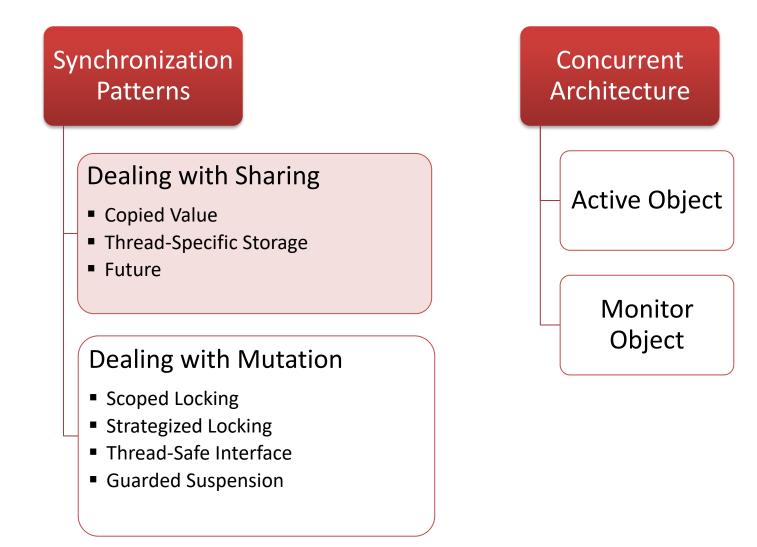
Concurrent Programming in Java<sup>\*\*</sup> Second Edition

**Design Principles and Patterns** 



Concurrent Programming in Java





#### **Copied Value**

There is no need to synchronize when a thread takes its arguments by copy and not by reference.

Data races or lifetime issues are not possible.

copiedValueDataRace.cpp

## **Thread-Specific Storage**

Thread-specific storage enables global state within a thread.

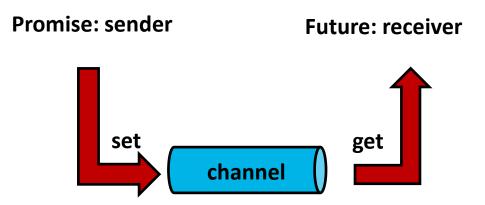
- Typical use-cases:
  - Porting a single-thread to multithreaded program
  - Compute thread-local and share the results
  - Thread-local logger

threadLocalSummation.cpp

#### Future

A future is a non-mutable placeholder for a value, which is set by a promise.

```
auto future = std::async([]{ return "LazyOrEager"; });
future.get();
```



#### Synchronization Patterns **Dealing with Sharing** Copied Value Thread-Specific Storage Future **Dealing with Mutation** Scoped Locking Strategized Locking

- Thread-Safe Interface
- Guarded Suspension

# Concurrent Architecture Active Object Monitor Object

# Scoped Locking

Scoped Locking is RAII applied to locking.

- Idea:
  - Bind the acquiring (constructor) and the releasing (destructor) of the resource to the lifetime of an object.
  - The lifetime of the object is bound.
  - The C++ run time is responsible for invoking the destructor and releasing the resource.
- C++ Implementation
  - std::lock\_guard and std::scoped\_lock
  - std::unique\_lock and std::shared\_lock

scopedLock.cpp

# Strategized Locking

#### Strategized Locking

- Enables it to use various locking strategies as replaceable components.
- Is the application of the strategy pattern to locking.
- Idea:
  - You want to use your library in various domains.
  - Depending on the domain, you want to use exclusive locking, shared locking, or no locking.
  - Configure your locking strategy at compile time or run time.

# Strategized Locking

Advantages:

Disadvantages:

- Run-time polymorphism
  - Enables it to change the locking strategy at runtime.
- Compile-time polymorphism
  - No cost at runtime
  - Flatter object hierarchies

- Run-time polymorphism
  - Needs a pointer indirection.

- Compile-time polymorphism
  - Produces in the error case a quite challenging to understand error message (when no concepts are used).

strategizedLockingRuntime.cpp
strategizedLockingCompileTimeWithConcepts.cpp

#### **Thread-Save Interface**

The thread-save interface extends the critical region to an object.

- Antipattern: Each member function uses a lock internally.
  - The performance of the system goes down.
  - Deadlocks appear when two member functions call each other.

#### **Thread-Save Interface**

#### A deadlock due to entangled calls.

```
struct Critical{
    void method1() {
        std::lock guard(mut);
        method2();
                                     int main() {
                                         Critical crit;
                                         crit.method1();
    }
    void method2() {
                                     }
        std::lock guard(mut);
    }
    std::mutex mut;
```

}

#### **Thread-Save Interface**

- Solutions:
  - All interface member functions (public) use a lock.
  - All implementation member functions (protected and private) must not use a lock.
  - The interface member functions call only implementation member functions.

#### **Guarded Suspension**

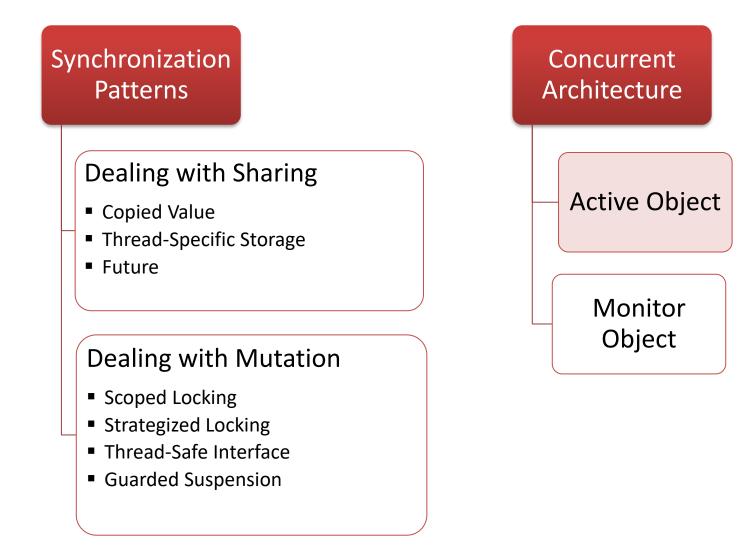
A guarded suspension consists of a lock and a condition. The condition has to be fulfilled by the calling thread.

- The calling thread will put itself to sleep if the condition is not meet.
- The calling thread uses a lock when it checks the condition.
- The lock protects the calling thread from a data race or deadlock.

### **Guarded Suspension**

- Guarded suspension enables thread synchronization. It is available in many variations.
  - The waiting thread is notified about the state change or asks for the state change.
    - Push principle: condition variables, future/promise pairs, atomics (C++20), or semaphores (C++20)
    - Pull principle: not natively supported in C++
  - The waiting thread waits with or without a time limit.
    - Condition variables, or future/promise pairs
  - The notification is sent to one or all waiting threads.
    - Shared futures, condition variables, atomics (C++20), or semaphores (C++20)

bossWorker.cpp



The active object pattern separates the method execution from the method invocation.

- Each object owns its own thread.
- Each method invocation is stored in an activation list.
- A scheduler triggers the method execution.

#### Proxy:

- Proxy for the member functions on the active object
- Triggers the construction of a request object which goes to the activation list and returns a future.
- It runs in the client thread.

#### Method Request

Includes all context information to be executed later.

Activation List:

- Has the pending requests objects.
- Decouples the client from the Active Object thread.

Scheduler:

- Runs in the thread of the Active Object.
- Decides with request from the Activation List is executes.

#### Servant:

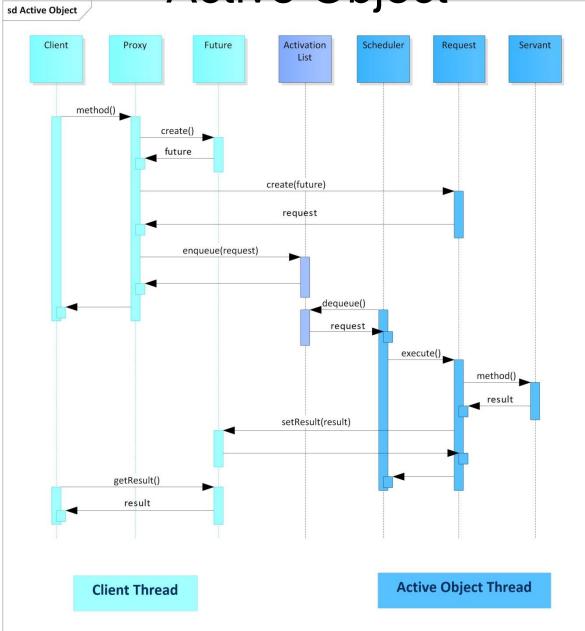
- Implements the member functions of the active objects.
- Supports the interface of the Proxy.

#### Future:

- Is created by the Proxy.
- Is only necessary if the request objects returns a result.
- The client uses the future to get the result of the request object.

#### **Dynamic Behavior**

- 1. Request construction and scheduling:
  - The client invokes the method on the proxy.
  - The proxy creates a request and passes it to the scheduler.
  - The scheduler enqueues the request on the activation list and returns a future to the client if the request returns something.
- 2. Member function execution
  - The scheduler determines which request becomes runnable.
  - It removes the request from the activation list and dispatches it to the servant.
- 3. Completion:
  - Stores eventually the result of the request object in the future.
  - Destructs the request object and the future if the client has the result.

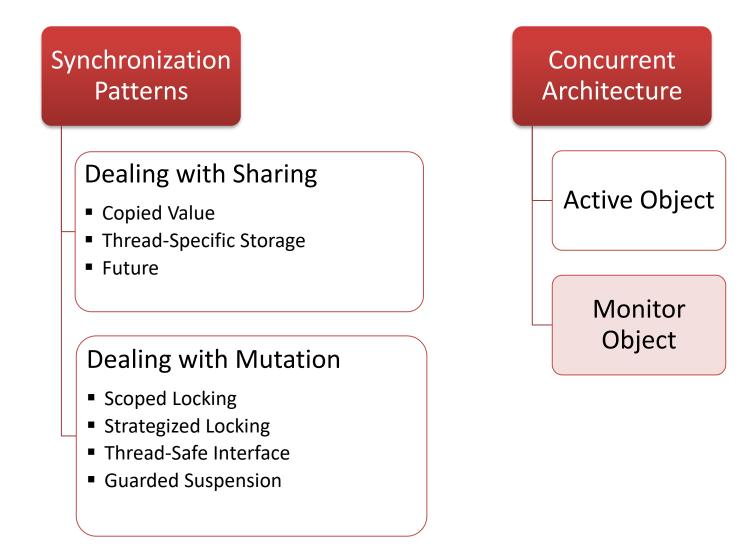


Advantages:

#### Disadvantages:

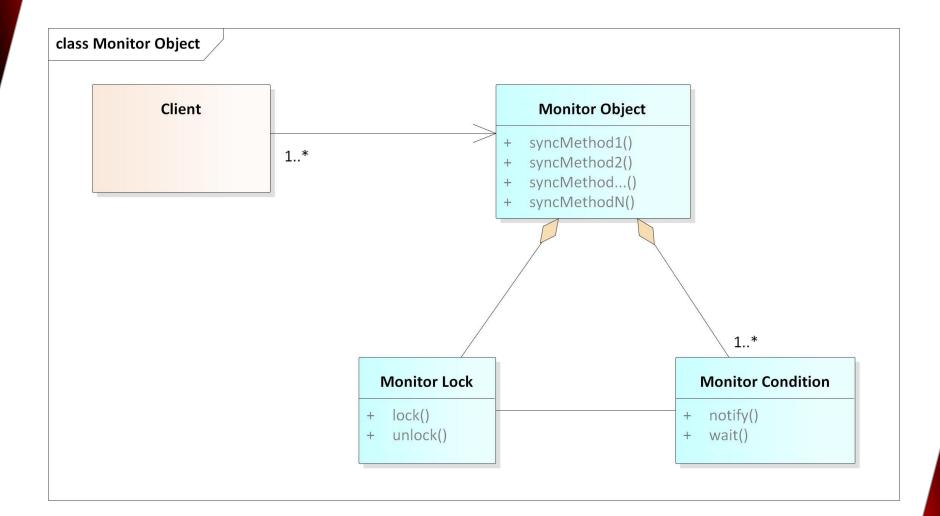
- Only the access to the activation list has to be synchronized
- Clear separation of client and server
- Improved throughput due to the asynchronous execution
- The scheduler can use various execution policies.

- If the member function execution is too fine-grained, the indirection may cause significant overhead.
- The asynchronous member function execution and the various execution strategies make the system quite difficult to debug.



The monitor object synchronizes the access to an object so that at most one member function can run at any moment in time.

- Each object has a monitor lock and a monitor condition.
- The monitor lock guarantees that only one client can execute a member function of the object.
- The monitor condition notifies the waiting clients.



Monitor Object:

• Support member functions, which can run in the thread of the client.

#### Synchronized Methods:

- Interface member functions of the monitor object.
- At most, one member function can run at any point in time
- The member functions should apply the thread-safe interface pattern.

Monitor Lock:

- Each monitor object has a monitor lock.
- Guarantees exclusive access to the member functions.

Monitor Condition:

- Allows various threads to store their member function invocation.
- When the current thread is done with its member function execution the next thread is awoken.

Advantages:

**Disadvantages:** 

- The synchronization is encapsulated in the implementation.
- The member function execution is automatically stored and performed.
- The monitor object is a simple scheduler.

- The synchronization mechanism and the functionality are strongly coupled and can, therefore, not easily be changed.
- When the synchronized member functions invoke an additional member function of the monitor object, a deadlock may happen.

monitorObject.cpp
monitorObjectCpp20.cpp

include automation

nt main(){

std::cout <<

std::vector

# www.ModernesCpp.com

std::function< bool(inti> myBindPr;

myVec.erase(std: nemove\_if(myVec.

std::cout << "myVde: ": for ( auto i: myVec)"std::cout << std::cout << "\n\n";</pre>

std::vector<int> myVec2(20); std::iota(myVec2/begin().tyVec2

std::cout << 'nyVec2: for ( auto in tyVec2) Rainer Grimm Training, Mentoring, and Technology Consulting <u>www.ModernesCpp.net</u>

nd( std::logical