

Best Practices for Concurrency

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Best Practices for Concurrency

General

Multithreading

Parallel

Memory Model

Best Practices for Concurrency

General

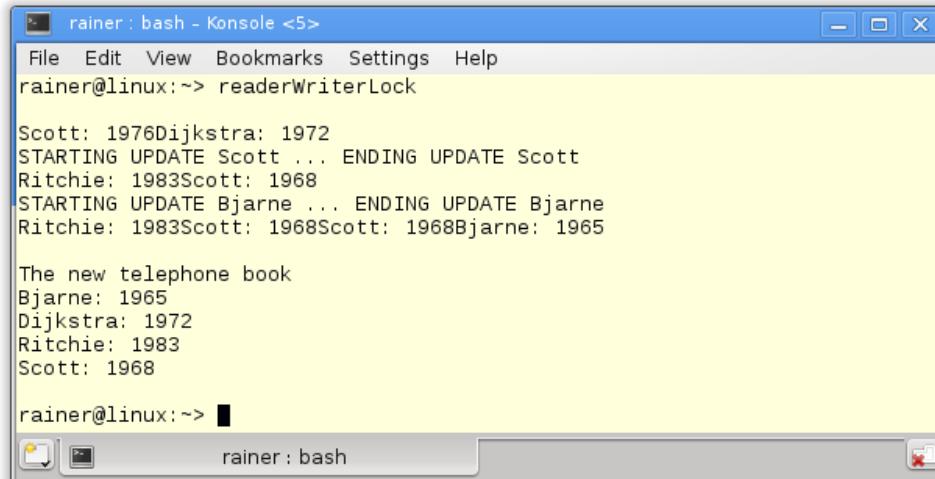
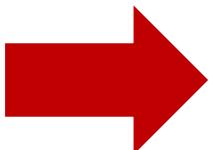
Multithreading

Memory Model

Code Reviews

```
map<string,int> teleBook{{"Dijkstra", 1972},  
                         {"Scott", 1976}, {"Ritchie", 1983}};  
  
shared_timed_mutex teleBookMutex;  
  
void addToTeleBook(const string& na, int tele){  
    lock_guard<shared_timed_mutex> writerLock(teleBookMutex);  
    cout << "\nSTARTING UPDATE " << na;  
    this_thread::sleep_for(chrono::milliseconds(500));  
    teleBook[na]= tele;  
    cout << " ... ENDING UPDATE " << na << endl;  
}  
  
void printNumber(const string& na){  
    shared_lock<shared_timed_mutex> readerLock(teleBookMutex);  
    cout << na << ":" << teleBook[na] << endl;  
}
```

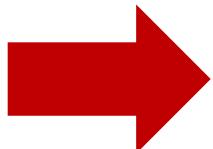
```
thread reader1([]{ printNumber("Scott"); });  
thread reader2([]{ printNumber("Ritchie"); });  
thread w1([]{ addToTeleBook("Scott",1968); });  
thread reader3([]{ printNumber("Dijkstra"); });  
thread reader4([]{ printNumber("Scott"); });  
thread w2([]{ addToTeleBook("Bjarne",1965); });  
thread reader5([]{ printNumber("Scott"); });  
thread reader6([]{ printNumber("Ritchie"); });  
thread reader7([]{ printNumber("Scott"); });  
thread reader8([]{ printNumber("Bjarne"); });  
  
reader1.join(), reader2.join();  
reader3.join(), reader4.join();  
reader5.join(), reader6.join();  
reader7.join(), reader8.join();  
w1.join(), w2.join();  
  
cout << "\nThe new telephone book" << endl;  
for (auto teleIt: teleBook){  
    cout << teleIt.first << ":" << teleIt.second << endl;  
}
```



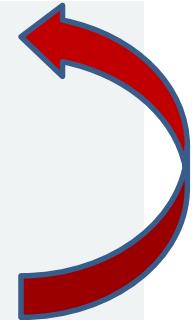
Code Reviews

```
map<string,int> teleBook{{"Dijkstra", 1972},  
                         {"Scott", 1976}, {"Ritchie", 1983}};  
  
shared_timed_mutex teleBookMutex;  
  
void addToTeleBook(const string& na, int tele){  
    lock_guard<shared_timed_mutex> writerLock(teleBookMutex);  
    cout << "\nSTARTING UPDATE " << na;  
    this_thread::sleep_for(chrono::milliseconds(500));  
    teleBook[na]= tele;  
    cout << " ... ENDING UPDATE " << na << endl;  
}  
  
void printNumber(const string& na){  
    shared_lock<shared_timed_mutex> readerLock(teleBookMutex);  
    cout << na << ":" << teleBook[na] << endl;  
}
```

```
thread reader1([]{ printNumber("Scott"); });  
thread reader2([]{ printNumber("Ritchie"); });  
thread w1([]{ addToTeleBook("Scott",1968); });  
thread reader3([]{ printNumber("Dijkstra"); });  
thread reader4([]{ printNumber("Scott"); });  
thread w2([]{ addToTeleBook("Bjarne",1965); });  
thread reader5([]{ printNumber("Scott"); });  
thread reader6([]{ printNumber("Ritchie"); });  
thread reader7([]{ printNumber("Scott"); });  
thread reader8([]{ printNumber("Bjarne"); });  
  
reader1.join(), reader2.join();  
reader3.join(), reader4.join();  
reader5.join(), reader6.join();  
reader7.join(), reader8.join();  
w1.join(), w2.join();  
  
cout << "\nThe new telephone book" << endl;  
for (auto teleIt: teleBook){  
    cout << teleIt.first << ":" << teleIt.second << endl;  
}
```



```
rainer@seminar:~> readerWriterLocks  
Bjarne: 0Ritchie: 1983  
STARTING UPDATE Scott ... ENDING UPDATE Scott  
STARTING UPDATE Bjarne ... ENDING UPDATE Bjarne  
Ritchie: 1983Scott: 1968Scott: 1968Scott: 1972Scott: 1968  
The new telephone book  
Bjarne: 1965  
Dijkstra: 1972  
Ritchie: 1983  
Scott: 1968  
rainer@seminar:~>
```



Minimise Sharing

- Summation of a vector with 100 000 000 elements

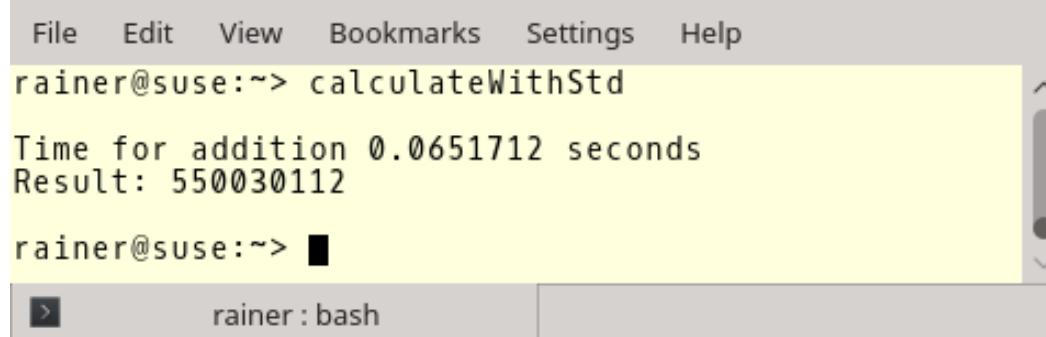
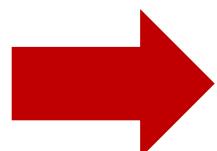
```
constexpr long long size = 100000000;  
...  
// random values  
std::vector<int> randValues;  
randValues.reserve(size);  
  
std::random_device seed;  
std::mt19937 engine(seed());  
std::uniform_int_distribution<> uniformDist(1, 10);  
for (long long i = 0 ; i < size ; ++i)  
    randValues.push_back(uniformDist(engine));  
  
...  
// calculate sum  
...
```

Minimise Sharing

- Single-threaded in two variations

```
unsigned long long sum {};
for (auto n: randValues) sum += n;

const unsigned long long sum = accumulate(randValues.begin(),
                                         randValues.end(), 0ll);
```



```
File Edit View Bookmarks Settings Help
rainer@suse:~> calculateWithStd
Time for addition 0.0651712 seconds
Result: 550030112
rainer@suse:~> █
▶ rainer:bash
```

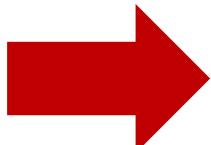
Minimise Sharing

- Four threads with a shared summation variable

```
void sumUp(unsigned long long& sum, const vector<int>& val,
           unsigned long long beg, unsigned long long end){
    for (auto it = beg; it < end; ++it){
        lock_guard<mutex> myLock(myMutex);
        sum += val[it];
    }
}

...
unsigned long long sum{};

thread t1(sumUp, ref(sum), ref(randValues), 0, fir);
thread t2(sumUp, ref(sum), ref(randValues), fir, sec);
thread t3(sumUp, ref(sum), ref(randValues), sec, thi);
thread t4(sumUp, ref(sum), ref(randValues), thi, fou);
```



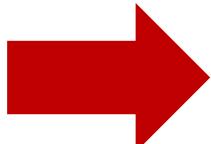
```
File Edit View Bookmarks Settings Help
rainer@suse:~> calculateWithLock
Time for addition 3.3389 seconds
Result: 549961505
rainer@suse:~> □
> rainer : bash
```

Minimise Sharing

- Four threads with a shared, atomic summation variable

```
void sumUp(atomic<unsigned long long>& sum, const vector<int>& val,
            unsigned long long beg, unsigned long long end){
    for (auto it = beg; it < end; ++it){
        sum += val[it];
    }
}

void sumUp(atomic<unsigned long long>& sum, const vector<int>& val,
            unsigned long long beg, unsigned long long end){
    for (auto it = beg; it < end; ++it){
        sum.fetch_add(val[it], memory_order_relaxed);
    }
}
```

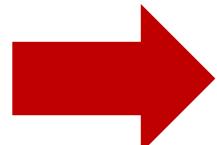


```
File Edit View Bookmarks Settings Help
rainer@suse:~> calculateWithAtomic
sum.is_lock_free(): true
Time for addition 1.33837 seconds
Result: 549992025
Time for addition 1.34625 seconds
Result: 549992025
rainer@suse:~> █
> rainer : bash
```

Minimise Sharing

- Four threads with a local summation variable

```
void sumUp(unsigned long long& sum, const vector<int>& val,
           unsigned long long beg, unsigned long long end){
    unsigned long long tmpSum{};
    for (auto i = beg; i < end; ++i){
        tmpSum += val[i];
    }
    lock_guard<mutex> lockGuard(myMutex);
    sum += tmpSum;
}
```



A screenshot of a terminal window titled "rainer@suse:~>". The window contains the following text:

```
File Edit View Bookmarks Settings Help
rainer@suse:~> localVariable
Time for addition 0.0284271 seconds
Result: 549996948
rainer@suse:~> █
```

The terminal window has a light gray background and a dark gray border. The text is in a monospaced font. A vertical scroll bar is visible on the right side of the window.

Minimise Sharing

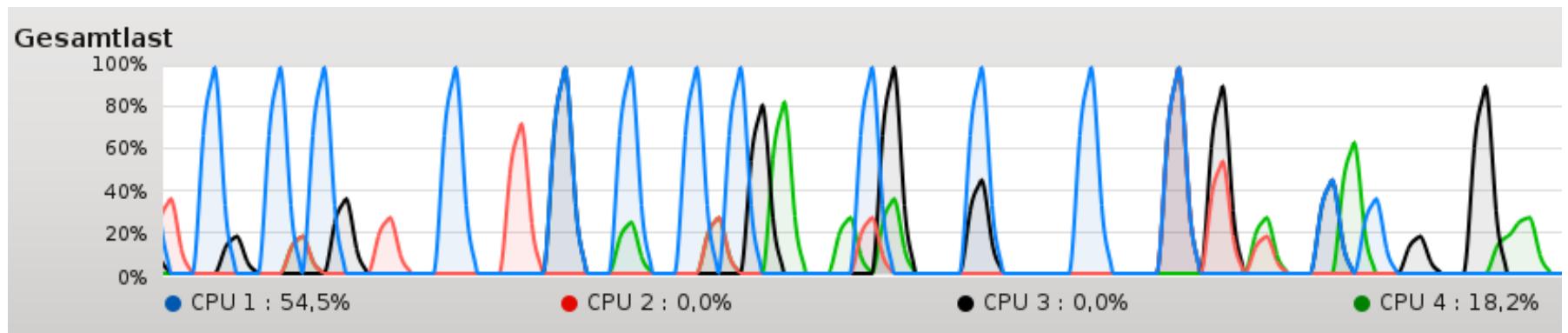
- Results

Single-threaded	4 threads with a lock	4 threads with an atomic	4 threads with a local variable
0.07 sec	3.34 sec	1.34 sec	0.03 sec

All Fine?

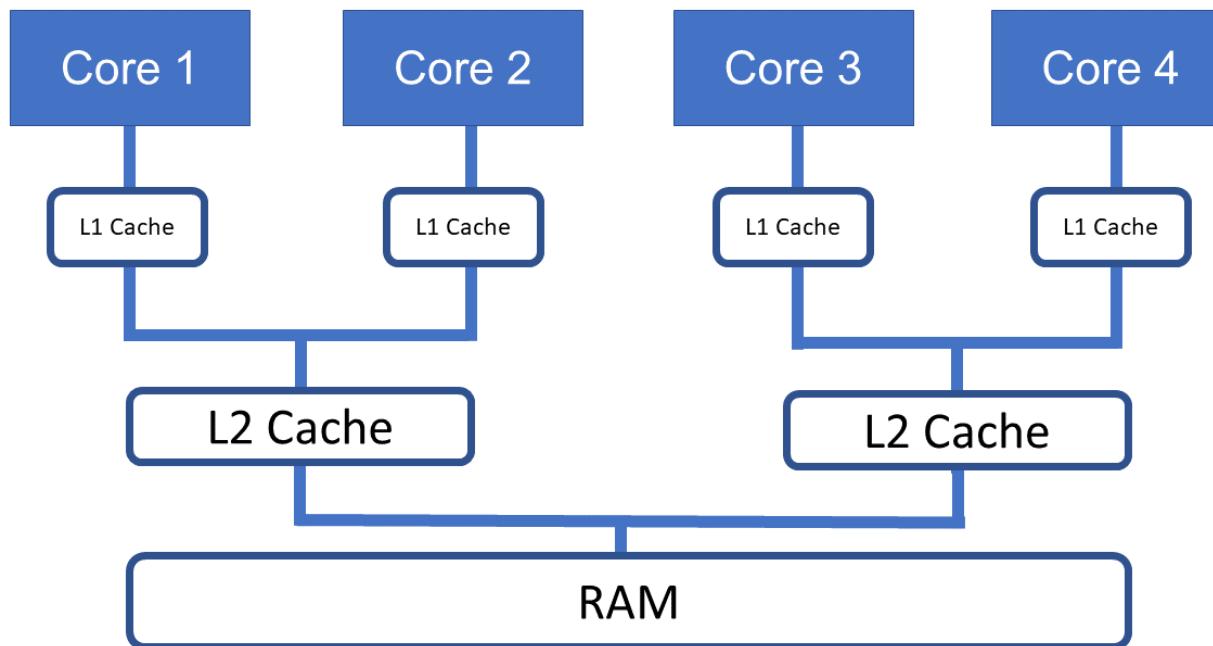
Minimise Sharing

- CPU Utilisation



Minimise Sharing

- Memory Wall



The RAM is the bottleneck.

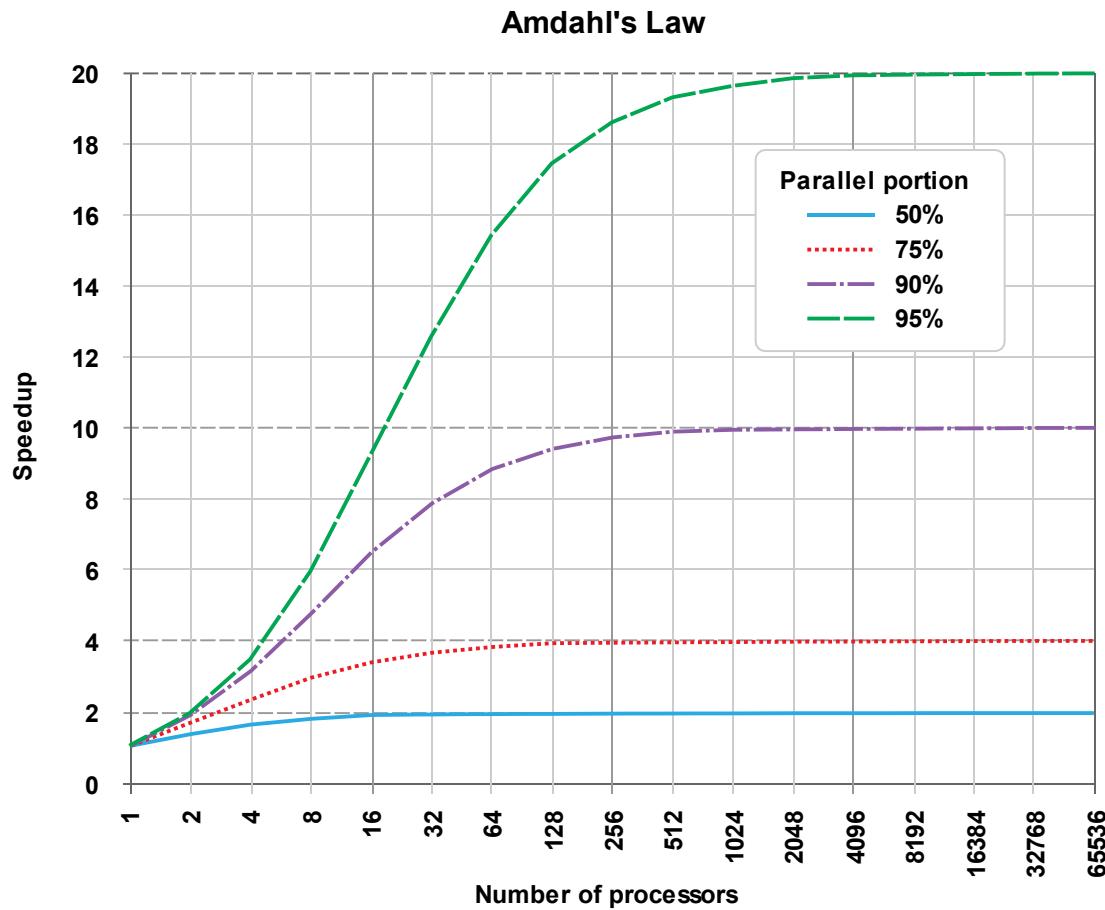
Minimise Waiting (Amdahl's Law)

$$\frac{1}{1 - p}$$

p: Parallel Code

$$p = 0.5 \rightarrow 2$$

Minimise Waiting (Amdahl's Law)



By Daniels220 at English Wikipedia, CC BY-SA 3.0, <https://commons.wikimedia.org/w/index.php?curid=6678551>

Use Code Analysis Tools (ThreadSanitizer)

Thread Sanitizer detects data races at runtime.

- Memory and performance overhead: 10x
- Available since GCC 4.8 and Clang 3.2

```
g++ conditionVariablePingPong.cpp -fsanitize=thread -g
-o conditionVariablePingPong -pthread
```

Use Code Analysis Tools (ThreadSanitizer)

```
bool dataReady= false;  
  
std::mutex mutex_;  
std::condition_variable condVar1;  
std::condition_variable condVar2;  
  
int counter=0;  
int COUNTLIMIT=50;  
  
void setTrue(){  
  
    while(counter <= COUNTLIMIT){  
  
        std::unique_lock<std::mutex> lck(mutex_);  
        condVar1.wait(lck,[]{return dataReady == false;});  
        dataReady= true;  
        ++counter;  
        std::cout << dataReady << std::endl;  
        condVar2.notify_one();  
  
    }  
}
```

```
void setFalse(){  
  
    while(counter < COUNTLIMIT){  
  
        std::unique_lock<std::mutex> lck(mutex_);  
        condVar2.wait(lck,[]{return dataReady == true;});  
        dataReady= false;  
        std::cout << dataReady << std::endl;  
        condVar1.notify_one();  
  
    }  
  
}  
  
int main(){  
  
    std::cout << std::boolalpha << std::endl;  
  
    std::cout << "Begin: " << dataReady << std::endl;  
  
    std::thread t1(setTrue);  
    std::thread t2(setFalse);  
  
    t1.join();  
    t2.join();  
  
    dataReady= false;  
    std::cout << "End: " << dataReady << std::endl;  
  
    std::cout << std::endl;  
}
```

Use Code Analysis Tools (ThreadSanitizer)

```
File Edit View Bookmarks Settings Help
rainer@linux: ~ conditionVariablePingPong
```

Begin: false
true

=====

WARNING: ThreadSanitizer: data race (pid=18139)

Read of size 4 at 0x000000004350 by thread T2:

#0 setFalse() /home/rainer/conditionVariablePingPong.cpp:30 (conditionVariablePingPong+0x000000401818)
#1 void std::__Bind_simple<void (*())()>::__M_invoke<(std::__Index_tuple<>) /usr/include/c++/6/functional:1400
#2 std::__Bind_simple<void (*())()>::operator()() /usr/include/c++/6/functional:1389 (conditionVariablePingPong+0x000000401818)
#3 std::thread::__State_impl<std::__Bind_simple<void (*())()>::__M_run() /usr/include/c++/6/thread:196 (conditionVariablePingPong+0x000000401818)
#4 <null> <null> (libstdc++.so.6+0x0000000c22de)

Previous write of size 4 at 0x000000604350 by thread T1 (mutexes: write M11):

#0 setTrue() /home/rainer/conditionVariablePingPong.cpp:21 (conditionVariablePingPong+0x00000040173d)
#1 void std::__Bind_simple<void (*())()>::__M_invoke<(std::__Index_tuple<>) /usr/include/c++/6/functional:1400
#2 std::__Bind_simple<void (*())()>::operator()() /usr/include/c++/6/functional:1389 (conditionVariablePingPong+0x00000040173d)
#3 std::thread::__State_impl<std::__Bind_simple<void (*())()>::__M_run() /usr/include/c++/6/thread:196 (conditionVariablePingPong+0x00000040173d)
#4 <null> <null> (libstdc++.so.6+0x0000000c22de)

Location is global 'counter' of size 4 at 0x000000604350 (conditionVariablePingPong+0x000000604350)

Mutex M11 (0x0000006042a0) created at:

#0 pthread_mutex_lock <null> (libtsan.so.0+0x00000003bc0f)
#1 __gthread_mutex_lock /usr/include/c++/6/x86_64-suse-linux/bits/gthr-default.h:748 (conditionVariablePingPong+0x000000401be0)
#2 std::mutex::lock() /usr/include/c++/6/bits/std_mutex.h:103 (conditionVariablePingPong+0x000000401be0)
#3 std::unique_lock<std::mutex>::lock() /usr/include/c++/6/bits/std_mutex.h:267 (conditionVariablePingPong+0x000000401be0)
#4 std::unique_lock<std::mutex>::unique_lock(std::mutex&) /usr/include/c++/6/bits/std_mutex.h:197 (conditionVariablePingPong+0x000000401be0)
#5 setTrue() /home/rainer/conditionVariablePingPong.cpp:18 (conditionVariablePingPong+0x0000004016f4)
#6 void std::__Bind_simple<void (*())()>::__M_invoke<(std::__Index_tuple<>) /usr/include/c++/6/functional:1400
#7 std::__Bind_simple<void (*())()>::operator()() /usr/include/c++/6/functional:1389 (conditionVariablePingPong+0x0000004016f4)
#8 std::thread::__State_impl<std::__Bind_simple<void (*())()>::__M_run() /usr/include/c++/6/thread:196 (conditionVariablePingPong+0x0000004016f4)
#9 <null> <null> (libstdc++.so.6+0x0000000c22de)

Thread T2 (tid=18140, running) created by main thread at:

#0 pthread_create <null> (libtsan.so.0+0x00000002b740)
#1 std::thread::__M_start_thread(std::unique_ptr<std::thread::__State, std::default_delete<std::thread::__State> 5d4)
#2 main /home/rainer/conditionVariablePingPong.cpp:49 (conditionVariablePingPong+0x00000040197c)

Thread T1 (tid=18139, running) created by main thread at:

#0 pthread_create <null> (libtsan.so.0+0x00000002b740)
#1 std::thread::__M_start_thread(std::unique_ptr<std::thread::__State, std::default_delete<std::thread::__State> 5d4)
#2 main /home/rainer/conditionVariablePingPong.cpp:48 (conditionVariablePingPong+0x00000040196b)

SUMMARY: ThreadSanitizer: data race /home/rainer/conditionVariablePingPong.cpp:30 in setFalse()

=====

false
true
false
true
false

```
rainer : bash
```

Use Code Analysis Tools (CppMem)

CppMem: Interactive C/C++ memory model

Model
standard preferred release_acquire tot relaxed_only

Program
examples/Paper | data_race.c

C Execution

```
// a data race (dr)
int main() {
    int x = 2;
    int y;
    {{ x = 3;
    ||| y = (x==3);
    }};
    return 0;
}
```

2 reset help **2 executions; 1 consistent, not race free**

Computed executions

Display Relations

sb asw dd cd
 rf mo sc lo
 hb vse ithb sw rs hrs dob cad
 unsequenced_races data_races

Display Layout

dot neato_par neato_par_init neato_downwards
 tex

[edit display options](#)

Execution candidate no. 2 of 2

3 previous consistent previous candidate next candidate next consistent 2 goto

Model Predicates

<input checked="" type="checkbox"/> consistent_race_free_execution = false	= true
<input checked="" type="checkbox"/> consistent_execution = true	
<input checked="" type="checkbox"/> assumptions	= true
<input checked="" type="checkbox"/> well_formed_threads	= true
<input checked="" type="checkbox"/> well_formed_rf	= true
<input checked="" type="checkbox"/> locks_only_consistent_locks	= true
<input checked="" type="checkbox"/> locks_only_consistent_lo	= true
<input checked="" type="checkbox"/> consistent_mo	= true
<input checked="" type="checkbox"/> sc_accesses_consistent_sc	= true
<input checked="" type="checkbox"/> sc_fenced_sc_fences_heeded	= true
<input checked="" type="checkbox"/> consistent_hb	= true
<input checked="" type="checkbox"/> consistent_rf	= true
<input checked="" type="checkbox"/> det_read	= true
<input checked="" type="checkbox"/> consistent_non_atomic_rf	= true
<input checked="" type="checkbox"/> consistent_atomic_rf	= true
<input checked="" type="checkbox"/> coherent_memory_use	= true
<input checked="" type="checkbox"/> rmw_atomicity	= true
<input checked="" type="checkbox"/> sc_accesses_sc_reads_restricted	= true
unsequenced_races	are absent
data_races	are present
ineterminate_reads	are absent
locks_only_bad_mutexes	are absent

4

```
graph TD; a[a:Wna x=2] -- sw --> b[b:Wna x=3]; b -- "rf,sw" --> c[c:Rna x=2]; c -- dr --> d[d:Wna y=0]; d -- sb --> a;
```

Files: [out.exc](#), [out.dot](#), [out.dsp](#), [out.tex](#)

Use Code Analysis Tools (CppMem)

```
int x = 0, std::atomic<int> y{0};

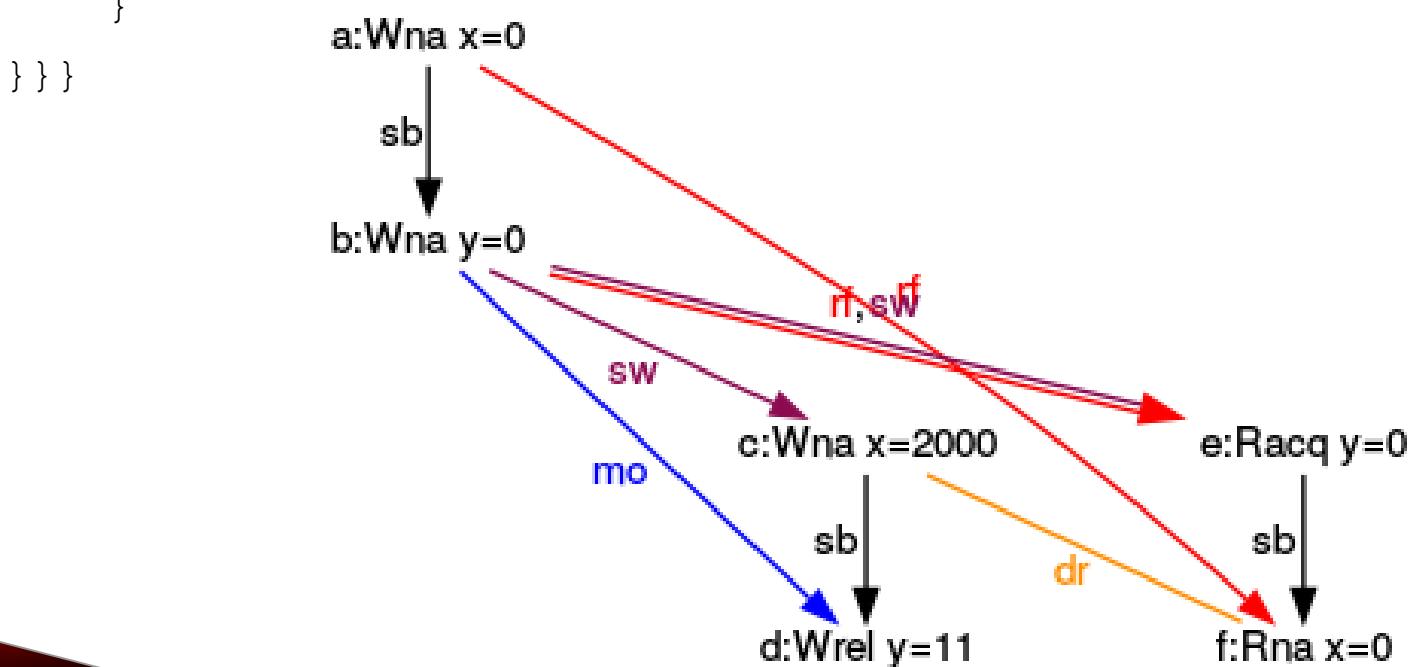
void writing() {
    x = 2000;
    y.store(11, std::memory_order_release);
}

void reading() {
    std::cout << y.load(std::memory_order_acquire) << " ";
    std::cout << x << std::endl;
}

std::thread thread1(writing);
std::thread thread2(reading);
thread1.join(), thread2.join();
```

Use Code Analysis Tools (CppMem)

```
int x = 0, atomic_int y = 0;  
{ {{ {  
    x = 2000;  
    y.store(11, memory_order_release);  
}  
||| {  
    y.load(memory_order_acquire);  
    x;  
}  
}}}
```



Use Immutable Data

Data Race: At least two threads access a shared variable at the same time. At least one thread tries to modify it.

Mutable?

Shared?

	No	Yes
No	OK	Ok
Yes	OK	Data Race

Use Immutable Data

Remaining Problem: initialise the data thread-safe

1. Early initialisation

```
const int val = 2011;  
thread t1([&val]{ .... });  
thread t2([&val]{ .... });
```

2. Constant expressions

```
constexpr auto doub = 5.1;
```

3. `call_once` and `once_flag`

```
void onlyOnceFunc() { .... };  
once_flag onceFlag;  
void func() { .... call_once(onceFlag,  
onlyOnceFunc); .... }  
thread t3{ func };  
thread t4{ func };
```

4. Static variables in a scope

```
void func() {  
.... static int val = 2011; ....  
}  
thread t5{ func() };  
thread t6{ func() };
```

Use Pure Functions

Pure Functions

Produce always the same result when given the same arguments

Have no side effect

Don't change the global state of the program

Added Value

- Easier to make correctness proofs
- Refactoring and testing is easier
- Results from previous function calls can be memorised
- **The function invocations can automatically be reordered or parallelised.**

Use Pure Functions

- Function

```
int powFunc(int m, int n){  
    if (n == 0) return 1;  
    return m * powFunc(m, n-1);  
}
```

- Metaprogramming

```
template<int m, int n>  
struct PowMeta{  
    static int const value = m * PowMeta<m, n-1>::value;  
};  
  
template<int m>  
struct PowMeta<m, 0>{  
    static int const value = 1;  
};
```

Use Pure Functions

- `constexpr` Function
 - almost pure functions

```
constexpr int powConst(int m, int n) {  
    int r = 1;  
    for(int k = 1; k <= n; ++k) r *= m;  
    return r;  
}
```

```
auto res = powConst(2, 10);  
auto res = PowMeta<2, 10>::value;  
constexpr auto res = powConst(2, 10);
```

Best Practices for Concurrency

General

Multithreading

Memory Model

Use Tasks instead of Threads

Threads

```
int res;  
thread t([&]{ res = 3 + 4; });  
t.join();  
cout << res << endl;
```

Tasks

```
auto fut = async([]{ return 3 + 4; });  
cout << fut.get() << endl;
```

Criteria	Thread	Task
Parties Involved	creator thread and child thread	promise and future
Communication	shared variable	communication channel
Thread Creation	obligatory	optional
Synchronisation	join call blocks	get call blocks
Exception in Child	creator thread and child thread die	return value of the promise
Forms of Communication	values	values, notifications, and exceptions

Use Tasks instead of Threads

C++20/23: Extend futures will support composition

- **then**: Execute the future if the previous future is done
- **when_any**: Execute the future if any of the previous future is done
- **when_all**: Execute the future if all of the previous futures are done

Don't use Fire and Forget Futures

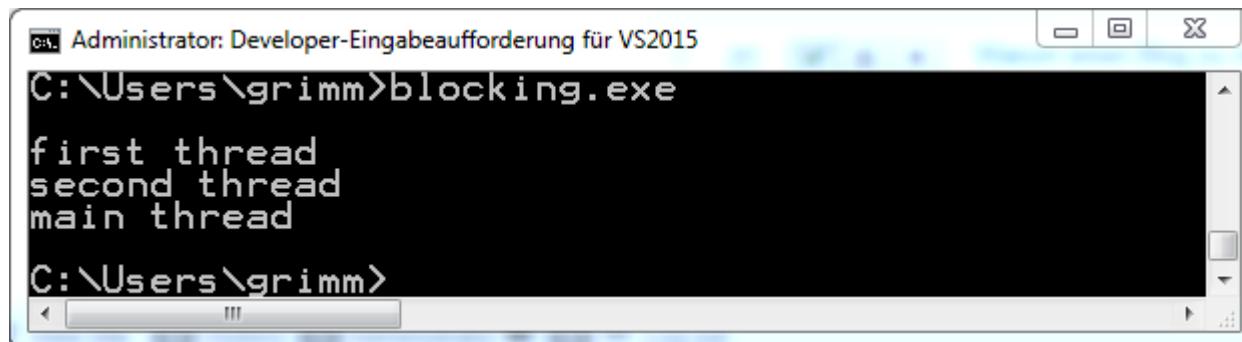
- The future created by `std::async` blocks until its associated promise is done.

```
#include <iostream>
#include <future>
int main() {
    std::cout << std::endl;
    std::async([]{ std::cout << "fire and forget" << std::endl; });
    std::cout << "main done " << std::endl;
}
```

➡ No `join` or `detach` call is required.
The asynchronous job is done synchronously.

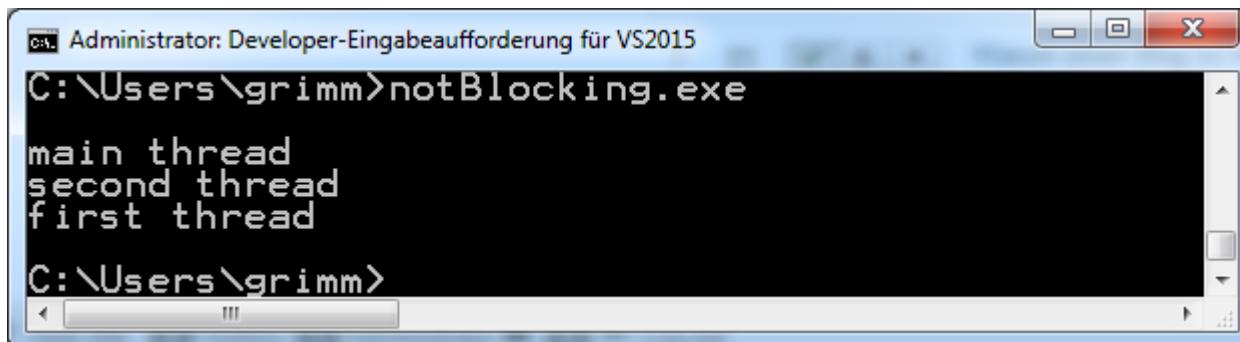
Don't use Fire and Forget Futures

```
int main() {  
    std::async(std::launch::async, [] {  
        std::this_thread::sleep_for(std::chrono::seconds(2));  
        std::cout << "first thread" << std::endl;  
    });  
    std::async(std::launch::async, [] {  
        std::this_thread::sleep_for(std::chrono::seconds(1));  
        std::cout << "second thread" << std::endl; }  
);  
    std::cout << "main thread" << std::endl;  
}
```



Don't use Fire and Forget Futures

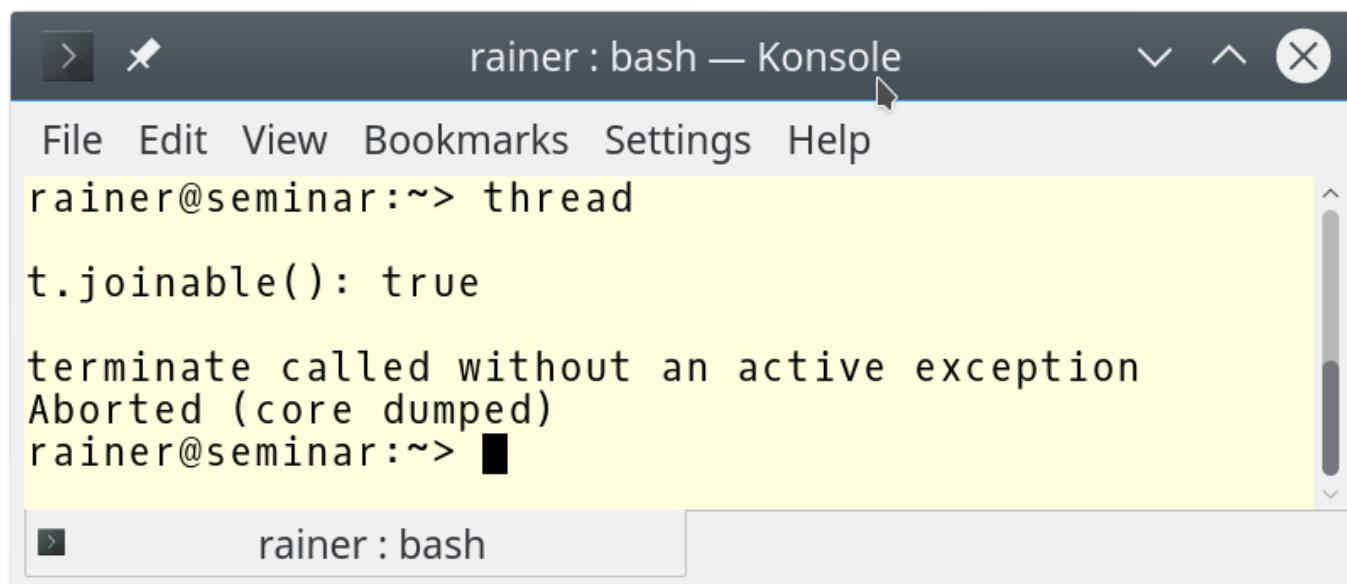
```
int main() {  
    auto first = std::async(std::launch::async, [] {  
        std::this_thread::sleep_for(std::chrono::seconds(2));  
        std::cout << "first thread" << std::endl;  
    });  
  
    auto second = std::async(std::launch::async, [] {  
        std::this_thread::sleep_for(std::chrono::seconds(1));  
        std::cout << "second thread" << std::endl;  
    });  
  
    std::cout << "main thread" << std::endl;  
}
```



std::jthread

Problem: std::thread **throws** std::terminate **in its destructor if still joinable.**

```
std::thread t[] { std::cout << "New thread"; } ;  
std::cout << "t.joinable(): " << t.joinable();
```



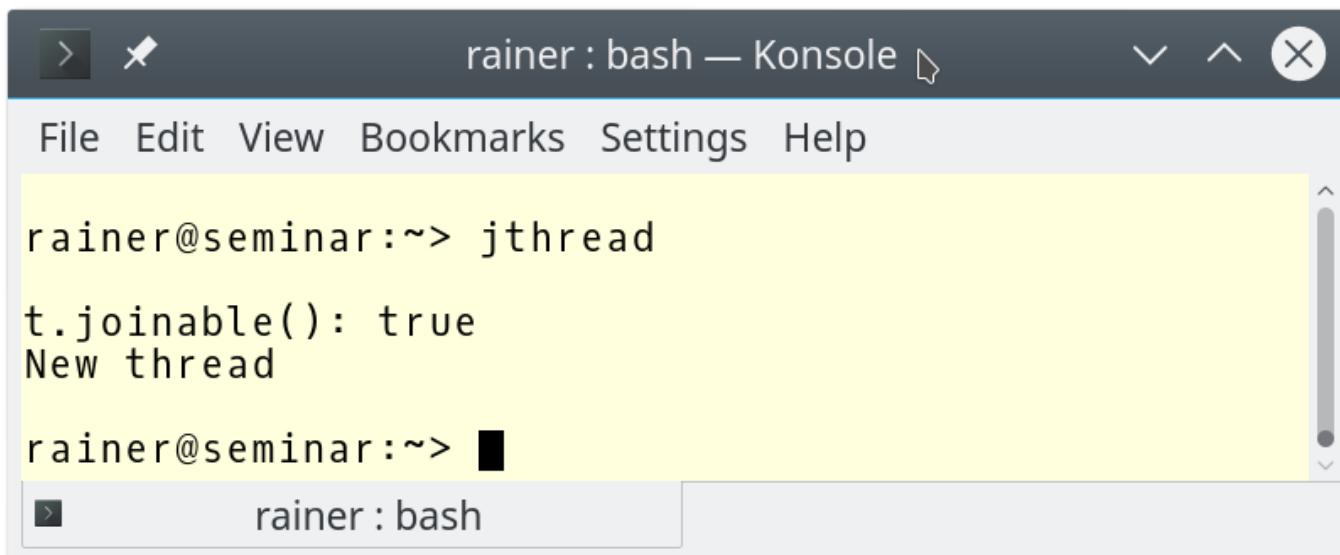
The screenshot shows a terminal window titled "rainer : bash — Konsole". The window has a menu bar with File, Edit, View, Bookmarks, Settings, and Help. The command "rainer@seminar:~> thread" is entered, followed by "t.joinable(): true". Then, an error message "terminate called without an active exception" and "Aborted (core dumped)" is displayed. The terminal window has a dark header and a light yellow body, with scroll bars on the right side.

```
rainer : bash — Konsole  
File Edit View Bookmarks Settings Help  
rainer@seminar:~> thread  
t.joinable(): true  
terminate called without an active exception  
Aborted (core dumped)  
rainer@seminar:~>
```

`std::jthread`

Solution: `std::jthread` (C++20) joins automatically at the end of its scope.

```
std::jthread t{[] { std::cout << "New thread"; } };
std::cout << "t.joinable(): " << t.joinable();
```



The screenshot shows a terminal window titled "rainer : bash — Konsole". The window contains the following text:

```
rainer@seminar:~> jthread
t.joinable(): true
New thread

rainer@seminar:~>
```

The terminal window has a dark header bar with icons for file operations and a close button. Below the header is a menu bar with "File", "Edit", "View", "Bookmarks", "Settings", and "Help". The main area of the terminal is light-colored with a vertical scroll bar on the right. The prompt "rainer@seminar:~>" appears twice, indicating the user's input and the program's output. The word "joinable()" is highlighted in blue, likely due to syntax highlighting in the terminal or the image itself.

Cooperative Interruption of std::jthread

- Instances of std::jthread can be interrupted

Receiver

- Explicit check:
 - `is_interrupted`: yields, when an interrupt was signalled
 - `wait` variations with predicate of `std::condition_variable`

Sender

- `interrupt`: signals an interrupt (and returns whether an interrupt was signaled before)

Cooperative Interruption of std::jthread

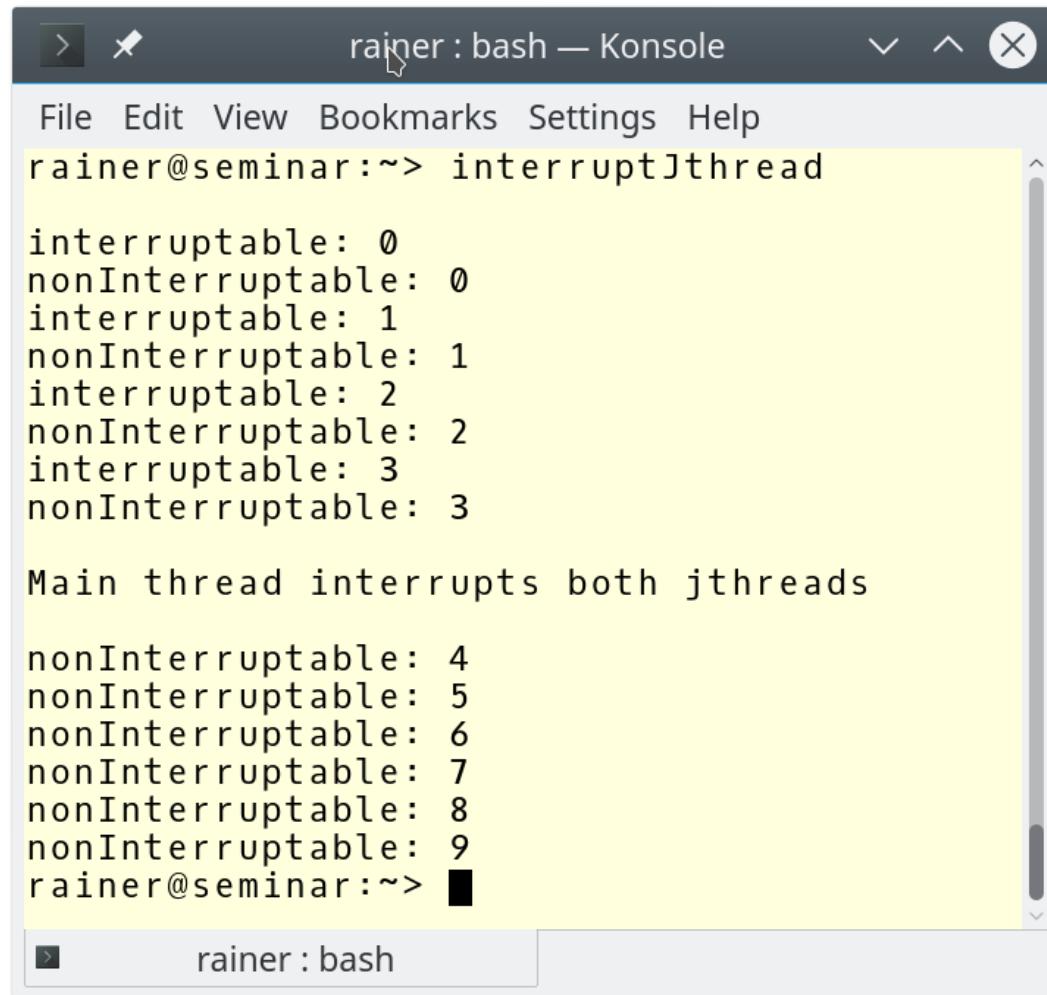
```
jthread nonInterruptable([]{
    int counter{0};
    while (counter < 10) {
        this_thread::sleep_for(0.2s);
        cerr << "nonInterruptable: "
             << counter << endl;
        ++counter;
    }
}) ;

jthread interruptable([](interrupt_token
    itoken) {
    int counter{0};
    while (counter < 10) {
        this_thread::sleep_for(0.2s);
        if (itoken.is_interrupted()) return;
        cerr << "interruptable: "
             << counter << endl;
        ++counter;
    }
}) ;

this_thread::sleep_for(1s);
cerr << endl;
cerr << "Main thread interrupts both jthreads" << endl;

nonInterruptable.interrupt();
interruptable.interrupt();
```

Cooperative Interruption of std::jthread



A screenshot of a terminal window titled "rainer : bash — Konsole". The window contains the following text output:

```
rainer@seminar:~> interruptJthread

interruptable: 0
nonInterruptable: 0
interruptable: 1
nonInterruptable: 1
interruptable: 2
nonInterruptable: 2
interruptable: 3
nonInterruptable: 3

Main thread interrupts both jthreads

nonInterruptable: 4
nonInterruptable: 5
nonInterruptable: 6
nonInterruptable: 7
nonInterruptable: 8
nonInterruptable: 9
rainer@seminar:> █
```

The terminal window has a yellow background for the main text area and a grey background for the command line prompt. A vertical scroll bar is visible on the right side of the yellow area.

Use Tasks instead of Condition Variables

Thread 1

```
{  
    lock_guard<mutex> lck(mut);  
    ready = true;  
}  
condVar.notify_one();
```

Thread 2

```
{  
    unique_lock<mutex>lck(mut);  
    condVar.wait(lck, []{ return ready; });  
}
```

prom.set_value();  fut.wait();

Criteria	Condition Variable	Tasks
Critical Region	Yes	No
Spurious Wakeup	Yes	No
Lost Wakeup	Yes	No
Repeatedly Synchronisation	Yes	No

Condition Variables are hard to use

Save Variant

```
{  
    lock_guard<mutex> lck(mut);  
    ready = true;  
}  
  
condVar.notify_one();  
  
...  
  
{  
    unique_lock<mutex>lck(mut);  
    condVar.wait(lck, []{ return ready; });  
}
```

Wrong Optimisation

```
std::atomic<bool> ready{false};  
  
...  
  
ready = true;  
condVar.notify_one();  
  
...  
  
{  
    unique_lock<mutex>lck(mut);  
    condVar.wait(lck, []{ return ready; });  
}
```

Condition Variables are hard to use

Wrong Optimisation

```
std::atomic<bool> ready{false};
```

...

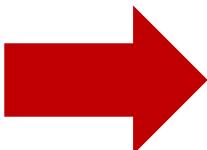
```
ready = true;  
condVar.notify_one();
```

Wrong Optimisation

```
std::unique_lock<std::mutex> lck(mut);  
while ( ![] { return dataReady.load(); } ()  
{  
    // time window  
    condVar.wait(lck);  
}
```



Condition variables are synchronisation mechanism at the **same** time.



Pack Mutexes into Locks

No Release of the lock

- std::mutex

```
mutex m;  
  
{  
    m.lock();  
    shaVar = getVar();  
    m.unlock();  
}
```

- std::lock_guard

```
mutex m;  
  
{  
    lock_guard<mutex> myLock(m);  
    shaVar = getVar();  
}
```

Minimal Locking

**Minimal locking or
Never call unknown code while holding a
lock**

- `std::lock_guard`

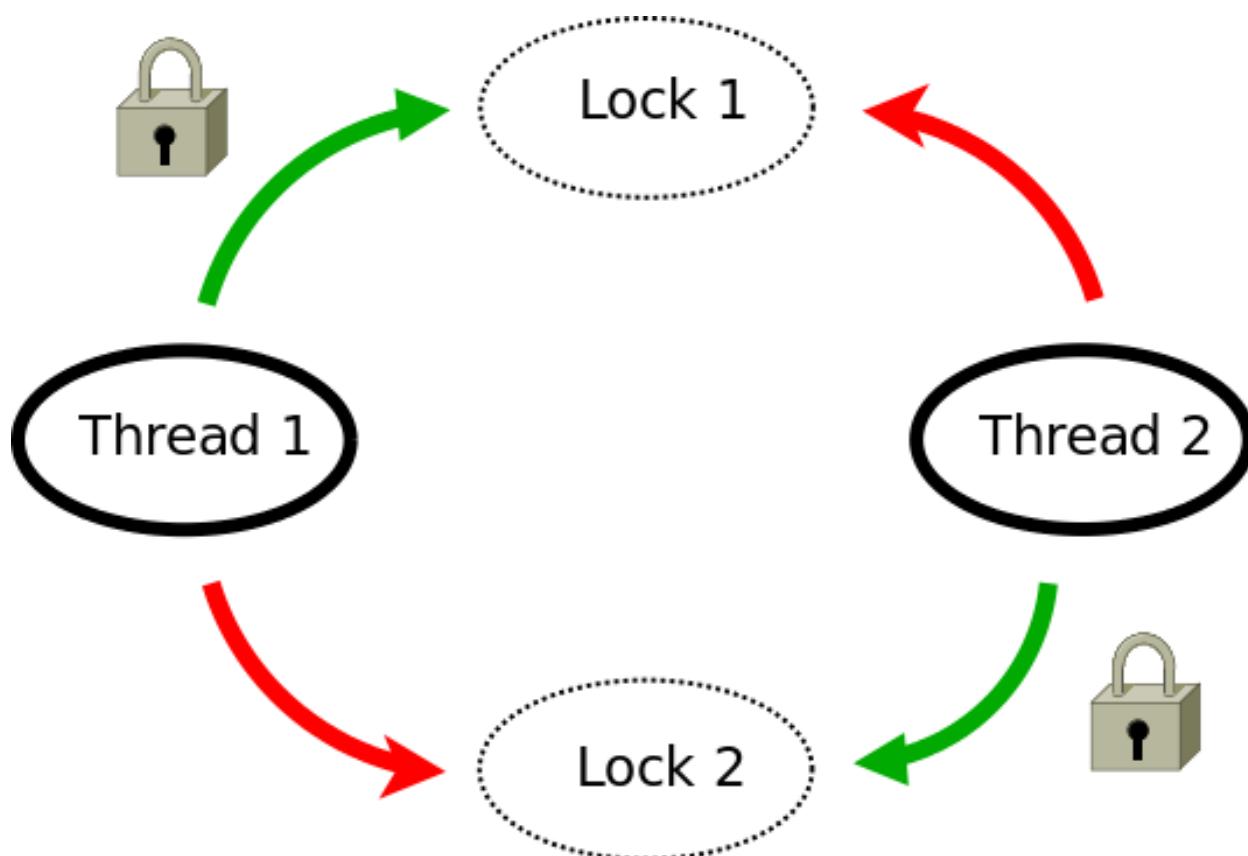
```
mutex m;  
{  
    lock_guard<mutex> myLock(m);  
    shaVar = getVar();  
}
```

- `std::lock_guard`

```
mutex m;  
auto localVar = getVar();  
{  
    lock_guard<mutex> myLock(m);  
    shaVar = localVar;  
}
```

Pack Mutexes into Locks

Locking of the mutexes is different order



Pack Mutexes into Locks

Atomic lock of the mutex

- std::unique_lock
 - {

 unique_lock<mutex> guard1(mut1, defer_lock);

 unique_lock<mutex> guard2(mut2, defer_lock);

 lock(guard1, guard2);

}
- std::scoped_lock (C++17)
 - {

 std::scoped_lock mut1, mut2;

}

Best Practices for Concurrency

General

Multithreading

Memory Model

Don't Program lock-free



Guide to Threaded Coding

1. Forget what you learned in Kindergarten
(ie stop Sharing)
2. Use Locks
3. Measure
4. Measure
5. Change your Algorithm
6. GOTO 1

∞. Lock-free

Lock-free coding is the last thing you want to do.

Red diagonal watermark "Tony Van Eerd" is overlaid on the right side.

Safety: off
How not to shoot yourself in the foot with C++ atomics

Anthony Williams



- Writing lock-free programs is hard **Fedor Pikus**
- Writing correct lock-free programs is even harder

The ugly side of weakly ordered atomics

Extreme complexity.

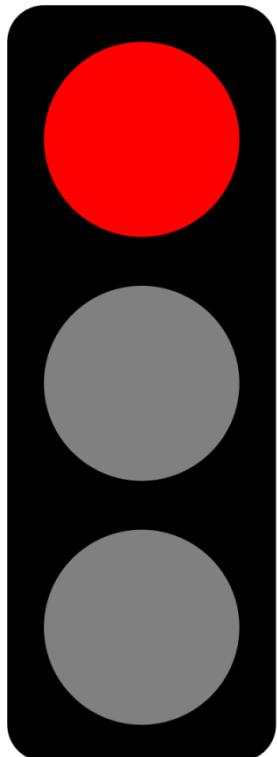
- The rules are not obvious.
- They're often downright surprising.
- And not even well understood.
- The committee still hasn't figured out how to define `memory_order_relaxed`.
... and I'm not even going to talk about `memory_order_consume`.

Hans Böhm

The specification of release-consume ordering is being revised, and the use of `memory_order_consume` is (since C++17) temporarily discouraged.

Don't Program lock-free: ABA

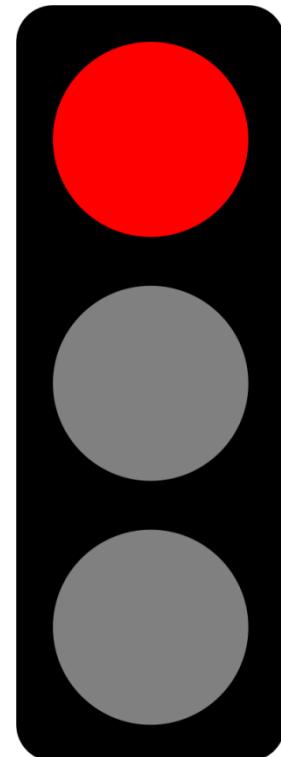
A



B



A



Don't Program lock-free: ABA

A lock-free, singly linked list (stack)



- Thread 1
 - wants to remove A
 - stores
 - head = A
 - next = B
 - checks if A == head
 - make B to the new head
 - B is already deleted by Thread 2
- Thread 2
 - removes A
 - removes B and deletes B
 - pushes A back

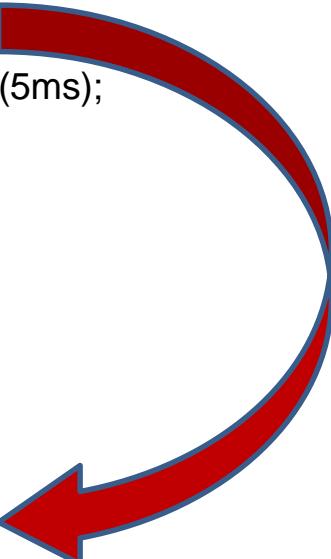
Use Proven Patterns

Wait with sequentiel consistency

```
std::vector<int> mySharedWork;
std::atomic<bool> dataReady(false);

void waitingForWork(){
    while ( !dataReady.load() ){
        std::this_thread::sleep_for(5ms);
    }
    mySharedWork[1] = 2;
}

void setDataReady(){
    mySharedWork = {1, 0, 3};
    dataReady.store(true);
}
```



```
int main(){
    std::thread t1(waitingForWork);
    std::thread t2(setDataReady);
    t1.join();
    t2.join();
    for (auto v: mySharedWork){
        std::cout << v << " ";           // 1 2 3
    }
};
```

Use Proven Patterns

Wait with acquire-release semantic

```
std::vector<int> mySharedWork;
std::atomic<bool> dataReady(false);

void waitForWork(){
    while ( !dataReady.load(std::memory_order_acquire) ){
        std::this_thread::sleep_for(5ms);
    }
    mySharedWork[1] = 2;
}

void setDataReady(){
    mySharedWork = {1, 0, 3};
    dataReady.store(true, std::memory_order_release);
}
```



```
int main(){
    std::thread t1(waitForWork);
    std::thread t2(setDataReady);
    t1.join();
    t2.join();
    for (auto v: mySharedWork){
        std::cout << v << " ";      // 1 2 3
    }
};
```

Use Proven Patterns

Atomic counter

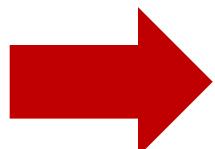
```
#include <vector>
#include <iostream>
#include <thread>
#include<atomic>

std::atomic<int> count{0};

void add() {
    for (int n = 0; n < 1000; ++n) {
        count.fetch_add(1, std::memory_order_relaxed);
    }
}

int main() {
    std::vector<std::thread> v;
    for (int n = 0; n < 10; ++n) {
        v.emplace_back(add);
    }
    for (auto& t : v) { t.join(); }
    std::cout << count;      // 10000
}
```

Don't Reinvent the Wheel



[Boost.Lockfree](#)
[CDS \(Concurrent Data Structures\)](#)

Don't Reinvent the Wheel

- Boost.Lockfree
 - Queue
 - A lock-free multi-producer/multi-consumer queue
 - Stack
 - A lock-free multi-producer/multi-consumer stack
 - spsc_queue
 - A wait-free single-producer/single-consumer queue (ringbuffer)

Don't Reinvent the Wheel

- Concurrent Data Structures (CDS)
 - Contains a lot of containers
 - Stacks (lock-free)
 - Queues and Priority-Queues (lock-free)
 - Ordered lists
 - Ordered sets and maps (lock-free and lock-based)
 - Unordered sets and maps (lock-free and lock-based)

Best Practices for Concurrency

General

Multithreading

Parallel

Memory Model

Blogs

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www.ModernesCpp.com [En]

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